Cybersecurity and Quantum Computation in Control of Cyberphysical Systems for Next-Generation Manufacturing

Helen Durand

Department of Chemical Engineering and Materials Science Wayne State University



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PAST EXPERIENCES

• B.S. Chemical Engineering, University of California, Los Angeles

• Materials and Processes Engineering Department, Aerojet Rocketdyne

• M.S. and Ph.D. Chemical Engineering, University of California, Los Angeles

• Assistant Professor, Wayne State University

INTRODUCTION

• Incentives for chemical process control



- Need for continuous monitoring and external intervention (process control)
- Objectives of a process control system
 - $\diamond\,$ Ensuring stability of the process
 - ♦ Suppressing the influence of external disturbances
 - ♦ Optimizing process performance



- How a feedback control loop (closed-loop system) works:
 - ♦ A variable describing the condition of a process (e.g., temperature, pressure, species concentration; known as an output) is measured by a sensor
 - The error between the measured output value and the desired value of this output (set-point) is calculated and fed to the controller
 - ◇ The controller computes a value of the manipulated input to the process to reduce the error
 - ◊ A control actuator (typically a valve) is used to apply the manipulated input value to the process



• Classical control: single-input/single-output (SISO) control design

- \diamond Proportional-integral-derivative (PID) control (error e(t))
 - \triangleright Error reflects difference between measured output and set-point
- \diamond Input/control action u(t)

$$u(t) = \underbrace{K_c e(t)}_{\mathbf{P}} + \underbrace{\frac{1}{\tau_I} \int_0^t e(\tau) d\tau}_{\mathbf{I}} + \underbrace{\tau_D \frac{de(t)}{dt}}_{\mathbf{D}}$$

 $\diamond K_c, \tau_I, \tau_D$: scalar values that can be picked (tuned)

ADVANCED MODEL-BASED PROCESS CONTROL

- Advanced process control utilizes a process dynamic model explicitly in the controller design
 - ♦ A mathematical process model is developed:
 - ▷ Constructed from first-principles
 - ▷ Identified from input-output process data
 - ◇ The model describes the process dynamics (variation of the process state variables in time due to disturbances, inputs, and interactions between variables)
 - ♦ Controllers are synthesized based on the process model
- Advantages of model-based control
 - ♦ Possibility of improved closed-loop performance
 - Model accounts for inherent process characteristics (e.g., nonlinear behavior, multivariable interactions)

NONLINEAR MODEL-BASED PROCESS CONTROL

• Example: continuous stirred tank reactor (CSTR)



• Model: system of nonlinear ordinary differential equations (ODEs)

$$\frac{dT}{dt} = \frac{F}{V_r}(T_0 - T) + \frac{(-\Delta H)}{\rho C_p} k_0 e^{-E/RT} C_A + \frac{Q}{\rho C_p V_r} \Rightarrow x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} T - T_s \\ C_A - C_{As} \end{bmatrix}, \ \dot{x} = \frac{dx}{dt}$$
$$\frac{dC_A}{dt} = \frac{F}{V_r}(C_{A0} - C_A) - k_0 e^{-E/RT} C_A \qquad \qquad u = Q - Q_s, \ w = C_{A0} - C_{A0s}$$

NONLINEAR MODEL-BASED PROCESS CONTROL

• Example: continuous stirred tank reactor (CSTR)



• Model: system of nonlinear ordinary differential equations (ODEs)

$$\dot{x} = f(x, u, w)$$

- Techniques for nonlinear controller design for driving the process state to the operating steady-state
 - \diamond Lyapunov-based control

 \diamond Model predictive control

NONLINEAR PROCESS SYSTEMS

• State-space description

$$\dot{x} = f(x, u, w)$$

♦ $x \in X \subset \mathbb{R}^n$ is the state, $u \in U \subset \mathbb{R}^m$ is the manipulated input, $w \in W \subset \mathbb{R}^l$ is the disturbance, f is a vector function

 Ω_{ρ}

- Explicit nonlinear feedback control law: u = h(x)
 - Control design technique: Lyapunov-based control (Y. Lin and E.D. Sontag, SCL, 1991; H. Khalil, Prentice Hall, 2002; P. D. Christofides and N. H. El-Farra, Springer-Verlag, 2005)
 - Renders the origin (steady-state) asymptotically stable
 - $\diamond\,$ There exists a Lyapunov function V which satisfies
 - $\dot{V} = \frac{\partial V(x)}{\partial x} f(x, h(x), 0) < 0, \ \forall x \in D$ V : energy of a physical system
 - ♦ Typically, $V(x) = x^T P x$ (quadratic) and $Ω_ρ ⊆ D$ is a level set of V where state constraints are met (i.e., $Ω_ρ := \{x : V(x) ≤ ρ\}$)
 - $\diamond \ u = h(x)$ possesses a degree of robustness to disturbances and uncertainty
- Performance considerations and constraints are not directly/explicitly taken into account

MODEL PREDICTIVE CONTROL



• Quadratic tracking stage cost:

 $l_T(x,u) = x^T Q x + u^T R u$

- $\diamond~Q,\,R$ are positive definite matrices
- Solve the optimization problem every Δ time units (sampling period)
 - $\diamond\,$ At each sampling time t_k



- Solution is a piecewise-constant input trajectory
 - \diamond Each piece is held constant for a period Δ
 - \diamond Prediction horizon N

MODEL PREDICTIVE CONTROL



solution

 \diamond Longer prediction horizon may

improve closed-loop performance

◇ Infinite/sufficiently long prediction

horizon

 \diamond Terminal cost/constraint

Steady-state

 t_{k+N}

Contractive constraint \diamond

NEXT-GENERATION MANUFACTURING

• Next-generation/smart manufacturing objectives (J. Davis, T. Edgar, J. Porter, J. Bernaden

and M. Sarli, Comput. Chem. Eng., 2012):

- \diamond Profitability
- \diamond Autonomy
- $\diamond\,$ Safety and cybersecurity



- Example: Moving away from a hierarchical approach to optimization and control
 - \diamond Upper layer:
 - Determine economicallyoptimal steady-state (real-time optimization (RTO)) (M. L. Darby,

M. Nikolaou, J. Jones and D. Nicholson, JPC,

2011)

- \diamond Lower layer:
 - Feedback control drives the state of the process to the optimal steady-state
- Tighter integration of plant operation and process economic optimization

PROCESS ECONOMICS AND CONTROL

• Traditional Paradigm



- Integration of economic optimization and process control
- Generalization of MPC
 - \diamond General (economic) stage cost



Steady-state operation

Dynamic/time-varying operation

- Economic MPC (EMPC) potential use cases:
 - ♦ Time-varying objective function or constraints (M. Ellis and P. D. Christofides, AIChE J.,

2013; A. Gopalakrishnan and L. T. Biegler, CACE, 2013)

(M. Ellis, H. Durand and P. D. Christofides, JPC, 2014)

ECONOMIC MPC FORMULATION

- EMPC formulation: $\min_{u(\cdot)\in S(\Delta)} \int_{t_k}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) d\tau$ s.t. $\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$ $\tilde{x}(t_k) = x(t_k)$ $u(t) \in U, \ \tilde{x}(t) \in X,$ $\forall t \in [t_k, t_{k+N})$ $|u(t_j) - u(t_{j-1})| \leq \epsilon_d$ $j = k, \dots, k+N-1$
- Components of EMPC:
 - \diamond Economic cost function
 - $\diamond\,$ Dynamic model
 - $\diamond\,$ State feedback measurement
 - $\diamond\,$ Input and state magnitude constraints
 - \diamond Input rate of change constraints
- System equipped with a measure of instantaneous economics l_e
- Computes control actions that optimize economics
- Accounts for input and state constraints
 - ♦ Examples: temperature or flow rate bounds
- Prevents rapid variations in inputs which may damage actuators

LYAPUNOV-BASED ECONOMIC MPC

Boundedness / Time-varying Operation (Mode 1) $\int_{t_{1}}^{t_{k+N}} l_{e}(\tilde{x}(\tau), u(\tau)) d\tau$ (M. Heidarinejad *et al.*, AIChE J., 2012) $\min_{u(\cdot)\in S(\Delta)}$ $\Omega_{
ho}$ $\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$ s.t. $x(t_k)$ $x(t_{k+1})$ $\tilde{x}(t_k) = x(t_k)$ $\Omega \rho_e$ $u(t) \in U, \ \tilde{x}(t) \in X, \ \forall \ t \in [t_k, t_{k+N})$ $|u_i(t_i) - h_i(\tilde{x}(t_i))| \le \epsilon_r, \ i = 1, \dots, m,$ $j = k, \ldots, k + N - 1$ $V(\tilde{x}(t)) \leq \rho_e, \ \forall \ t \in [t_k, t_{k+N})$ if $V(x(t_k)) \leq \rho_e$ and $t_k < t_s$

- Provable stability: boundedness of the closed-loop state in Ω_{ρ} ($\Omega_{\rho_e} \subset \Omega_{\rho}$)
- Provable feasibility: h(x) meets all state and input constraints

LYAPUNOV-BASED ECONOMIC MPC

Convergence to the Steady-State (Mode 2) $\min_{u(\cdot)\in S(\Delta)} \quad \int_{t_k}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) \ d\tau$ $\Omega_{
ho}$ $\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$ s.t. $x(t_k)$ $\tilde{x}(t_k) = x(t_k)$ $x(t_{k+1})$ $u(t) \in U, \ \tilde{x}(t) \in X, \ \forall \ t \in [t_k, t_{k+N})$ $|u_i(t_j) - h_i(\tilde{x}(t_j))| \le \epsilon_r, \ i = 1, \dots, m,$ $j = k, \ldots, k + N - 1$ $\frac{\partial V(x(t_k))}{\partial r}f(x(t_k), u(t_k), 0)$ $\leq \frac{\partial V(x(t_k))}{\partial x} f(x(t_k), h(x(t_k)), 0)$ if $V(x(t_k)) > \rho_e$ or $t_k \ge t_s$

- Compute control actions that decrease the Lyapunov function
- Provable stability: convergence to a small neighborhood of the steady-state

CYBERSECURITY AND PROCESS CONTROL SYSTEMS

- Cyberattacks on control systems seek to impact a physical process and can impact safety, profit, and production rates (A.A. Cárdenas *et al.*, *ASIACCS*, 2011)
- Do cyberattackers care about attacking control and manufacturing systems?
 - \diamond 2010: Stuxnet (trellix.com)
 - \triangleright Attack on Iranian nuclear facilities
 - $\triangleright\,$ Worm entered systems via USB sticks and spread
 - \triangleright Searched for control system software
 - \triangleright Ran centrifuges at conditions that cause breakdown
 - Falsified information to main controller so that there was no indication of a problem
 - ♦ December 2015: Part of Ukraine power grid (K. Zetter, Wired, 2016)
 - ▷ Remote manipulation of circuit breakers
 - ▷ Locking real operators out of their accounts
 - ▷ Malicious firmware prevented operators from un-doing attacks
 - ▷ Turned off backup power for operators
 - > Telephone denial of service to prevent operators from finding out about power outages too quickly

CYBERSECURITY AND PROCESS CONTROL SYSTEMS

- Cyberattacks on control systems seek to impact a physical process and can impact safety, profit, and production rates (A.A. Cárdenas et al., ASIACCS, 2011)
- Do cyberattackers care about attacking control and manufacturing systems?
 - ♦ 2017: Triton (M. Giles, MIT Technology Review, 2019)
 - Malware that can prevent safety instrumented systems from activating when needed
 - \triangleright Present on a petrochemical plant in Saudi Arabia
 - Flaw caused safety systems to act up in a way that revealed it before it could cause an incident
 - \diamond 2021: Florida water treatment plant (J. Bergal, PEW, 2021)
 - Remote user changed sodium hydroxide level to be 100 times higher than it should have been
 - \triangleright Operator saw this and changed it back
 - ♦ 2021: Colonial Pipeline (W. Turton and K. Mehrotra, Bloomberg, 2021)
 - ▷ Ransom note requesting payment appeared on company computer
 - Company closed down pipeline due to uncertainty as to whether operational technology was compromised

CYBERATTACK-RESILIENT CHEMICAL PROCESSES

- Examples of attack types: (N. Tuptuk and S. Hailes, Journal of Manufacturing Systems, 2018)
 - \diamond Denial of Service: Preventing parts of a network from delivering to others

 - \diamond False data injection
 - ♦ Time delay attack: Delay occurs in measurements or control actions
 - ♦ Data tampering attack: Data can be altered in storage or transmittal
 - ◇ Replay attack: Correct information from before is sent again
- Cyberattacks on feedback controllers are problematic because they remove associations between state measurements and inputs
 - \diamond Undesired inputs $u \in U$ can be applied at a given state
 - $\diamond\,$ Defies standard notions of feedback control
- Desirable to understand how elements of a control loop can contribute to detection and handling of attacks
 - ♦ Goal: Understand how and whether control theory-based cyberattack-handling can aid in providing security with flexibility for next-generation manufacturing

CYBERATTACK-RESILIENT CHEMICAL PROCESSES: A NONLINEAR SYSTEMS DEFINITION

(H. Durand, Mathematics, 2018)

- Physical damage from attacks can come from manipulating actuators in a rogue manner (directly or indirectly)
- Focus on sensor and actuator attacks individually to build toward handling both at once
- Cyberattack-resilience for state measurement falsification requires:
 - ♦ There exist no possible input policies given the controllers used and their implementation strategies such that $x(t) \notin X$, for any allowable initial state $x_0 \in \overline{X}$ and $w(t) \in W$, $t \in [0, \infty)$



DISCOVERING PROPERTIES OF CYBERATTACK-RESILIENT PROCESS CONTROL DESIGNS

- The definition of cyberattack-resilient control design is non-constructive
- Developing cyberattack-resilient control strategies will require a better understanding of which designs do and do not work and why
- Explore 2 ideas for cyberattack-resilient controllers:
 - ♦ Controller implementation incorporating randomness
 - \diamond Integrating feedback control/open-loop control
- Conclusions:
 - ♦ Nonlinear systems definition of cyberattack-resilience must be met
 - \triangleright Hoping the attacker lacks knowledge about the control design is insufficient
 - ◊ Other techniques (e.g., process design perspectives or techniques which combine control with detection) should be investigated

CONTROLLER IMPLEMENTATION INCORPORATING RANDOMNESS

- Attacks may be designed by reverse engineering known control laws
 - ♦ Suggests that randomly selecting the controller to be used at a given sampling time may make cyberattack design more difficult
 - ♦ Randomness in control design can only be considered if closed-loop stability is maintained under normal operation
 - Closed-loop stability and feasibility guarantees can be made with a randomized LEMPC implementation strategy
 - \triangleright Cyberattack-resiliency is not guaranteed



- Implementation strategy:
 - \diamond Develop n_p LEMPC's and $h_1(x)$
 - \diamond At each t_k , randomly select one of the controllers until one is found for which:
 - $\triangleright x(t_k) \in \Omega_{\rho_i}, i = 1, \dots, n_p, \text{ for the } n_p th$ LEMPC
 - $\triangleright x(t_k) \in \Omega_{\rho_1} \text{ for } h_1(x)$

Process Description

• Continuous stirred tank reactor (CSTR) with second-order, exothermic, irreversible reaction of the form $A \rightarrow B$:

$$\frac{dC_A}{dt} = \frac{F}{V}(C_{A0} - C_A) - k_0 e^{\frac{-E}{RT}} C_A^2$$
$$\frac{dT}{dt} = \frac{F}{V}(T_0 - T) + \frac{-\Delta H}{\rho_L C_p} k_0 e^{\frac{-E}{RT}} C_A^2 + \frac{Q}{\rho_L C_p V}$$

• Control objective: regulate the process in an economically optimal time-varying fashion while maintaining closed-loop stability

 \diamond Economic cost:

$$\int_{t_k}^{t_{k+N}} \left[k_0 e^{-\frac{E}{RT(\tau)}} C_A(\tau)^2\right] d\tau$$

 \diamond Manipulated input constraints

$$0.5 \le C_{A0} \le 7.5 \; {
m kmol/m^3} \; -5.0 imes 10^5 \le Q \le 5.0 imes 10^5 \; {
m kJ} \; / \; {
m h}$$

♦ Deviation variables:

$$x_1 = C_A - C_{As}, \quad x_2 = T - T_s$$

 \diamond Process model in input-affine form $\dot{x} = \tilde{f}(x) + gu$

Lyapunov-Based Controller Design

- Lyapunov-based controller for the inlet concentration: $h_{1,1}(x) = 0 \text{ kmol/m}^3$
 - ♦ Lyapunov-based controller for the heat rate input:
 - ▷ Sontag's Formula (Y. Lin and E.D. Sontag, SCL, 1991)

$$h_{2,1}(x) = \begin{cases} -\frac{L_{\tilde{f}}V_1 + \sqrt{L_{\tilde{f}}V_1^2 + L_{g_2}V_1^4}}{L_{g_2}V_1}, & \text{if } L_{g_2}V_1 \neq 0\\ 0, & \text{if } L_{g_2}V_1 = 0 \end{cases}$$

♦ A quadratic Lyapunov function of the form $V_1(x) = x^T P x$ with:

$$P = \left[\begin{array}{rrr} 1200 & 5\\ 5 & 0.1 \end{array} \right]$$

♦ Stability region $ρ_1 = 180$ (i.e., $Ω_{ρ_1} = \{x \in R^2 : V_1(x) \le ρ_1\}$)

- Process state initialized at $x_{init} = [-0.4 \text{ kmol/m}^3 \text{ } 20 \text{ K}]^T$
- LEMPC parameters: $N = 10, \Delta = 0.01$ h
- Process simulated with an integration step size of 10^{-4} h

Randomized LEMPC Development

- 6 LEMPC's were designed
 - $\diamond \ \Omega_{\rho_i} \subseteq \Omega_{\rho_1}, \ i = 1, \dots, 6$
 - $\diamond \ h_{i,1} = 0 \ \rm kmol/m^3$
 - $\begin{tabular}{ll} \diamond $h_{i,2}$ designed via \\ Sontag's control law \end{tabular} \end{tabular}$



Randomized LEMPC and LEMPC Under a Cyberattack

- Cyberattack with $x_f = [-0.0521 \text{ kmol/m}^3 8.3934 \text{ K}]^T$ is applied to a single LEMPC and the randomized LEMPC implementation strategy
- Randomized LEMPC results depend on seed to random number generator
- Randomized LEMPC barely delayed the time until $x_2 > 55$ K compared to the single LEMPC (0.0142 h)



Seed	Time $x_2 > 55$ (h)
5	0.0231
10	0.0144
15	0.0142
20	0.0323
25	0.0247
30	0.0142
35	0.0142
40	0.0146
45	0.0247
50	0.0142

INTEGRATING FEEDBACK CONTROL/OPEN-LOOP CONTROL

- Randomized LEMPC implementation strategy could not guarantee that no problematic inputs could be applied over time (even for steady-state tracking)
- Cyberattack resilience against state measurement falsification could be achieved for systems with an open-loop stable steady-state
 - \diamond Applying the steady-state input u_s by passes the issues with cyberattacks on feedback and drives the closed-loop state to the origin
 - $\diamond\,$ Loses benefits of feedback control
- Cyberattack-resilience definition must be met
- Three concepts for utilizing LEMPC to attempt to detect attacks were explored (H. Durand and M. Wegener, *Mathematics*, 2020; H. Oyama and H. Durand, *AIChE J.*, 2020)
 - ♦ LEMPC with random control law modifications to probe for cyberattacks
 - State feedback LEMPC with an attack detection strategy based on state predictions at each sampling time
 - Output feedback LEMPC (M. Ellis, J. Zhang, J. Liu and P. D. Christofides, SCL, 2014; L. Lao, M. Ellis, H. Durand and P. D. Christofides, AIChE J., 2015) with an attack detection strategy based on redundant state estimators

OBSERVABILITY ASSUMPTION

• M sets of measurements are continuously available:

$$y_i(t) = k_i(x(t)) + v_i(t)$$

- $\diamond k_i$ is vector-valued function, and v_i represents the measurement noise associated with the measurements y_i
- $\diamond v_i \in V_i \subset \mathbb{R}_i^q \ (|v_i| \le \theta_{v,i}), \ i = 1, \dots, M$
- A deterministic observer exists for each of the *M* sets of measurements:

 $\dot{z}_i = F_i(\epsilon_i, z_i, y_i)$

- \diamond Observer estimate z_i ; $\epsilon_i > 0$
- Assumptions:
 - \diamond For an initial state estimate with sufficiently low error between z_i and x, $h(z_i)$ maintains the closed-loop state in Ω_{ρ}
 - \diamond There exists a time t_{bi} such that:

 $|z_i(t) - x(t)| \le \epsilon_{mi}$



CYBERATTACK-RESILIENT OUTPUT FEEDBACK LEMPC

- Cyberattacks on state measurements could impact the state estimate used by the LEMPC
- If the estimate is sufficiently incorrect, the closed-loop state may exit Ω_{ρ}
- Estimator properties suggest an attack detection methodology

$$\diamond |z_i(t) - x(t)| \leq \max\{e_{mi}\}, \quad i = 1, \dots, M$$

- ♦ Implies $|z_i(t) z_j(t)| \le \epsilon_{\max}$, $i, j = 1, \ldots, M$, when no attack occurs
- Condition can be used with redundant estimators to attempt to flag falsified sensor measurements

 $\int_{t}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) \, d\tau$ $\min_{\substack{u(t)\in S(\Delta)}}$ s.t. $\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t))$ $\tilde{x}(t_k) = z_1(t_k)$ $\tilde{x}(t) \in X, \forall t \in [t_k, t_{k+N})$ $u(t) \in U, \ \forall t \in [t_k, t_{k+N})$ $V(\tilde{x}(t)) \le \rho_{e,1}, \quad \forall t \in [t_k, t_{k+N}),$ if $\tilde{x}(t_k) \in \Omega_{\rho_{e,1}}$ $\frac{\partial V(\tilde{x}(t_k))}{\partial x}(f(\tilde{x}(t_k), u(t_k)))$ $\leq \frac{\partial V(\tilde{x}(t_k))}{\partial x} \left(f(\tilde{x}(t_k), h(x(t_k))) \right)$ if $\tilde{x}(t_k) \in \Omega_{\rho} / \Omega_{\rho_{e,1}}$

CYBERATTACK-RESILIENT OUTPUT FEEDBACK LEMPC

- Consider that at least one state estimate is not impacted by an attacker
- If $|z_i(t) z_j(t)| > \epsilon_{\max}$, $i, j = 1, \dots, M$, flag an attack
- If $|z_i(t) z_j(t)| \le \epsilon_{\max}, i, j = 1, \dots, M$, but an attack occurred:
 - $\label{eq:closed-loop state will be maintained} $$ $$ $$ Closed-loop state will be maintained $$ $$ in Ω_{ρ} over the subsequent sampling $$ period under sufficient conditions $$ $$ $$
 - \triangleright Examples: sufficiently small $\rho_{e,1}$, θ , and Δ

 $\min_{u(t)\in S(\Delta)} \int_{t_{i}}^{t_{k+N}} l_{e}(\tilde{x}(\tau), u(\tau)) d\tau$ s.t. $\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t))$ $\tilde{x}(t_k) = z_1(t_k)$ $\tilde{x}(t) \in X, \forall t \in [t_k, t_{k+N})$ $u(t) \in U, \ \forall t \in [t_k, t_{k+N})$ $V(\tilde{x}(t)) \le \rho_{e,1}, \quad \forall t \in [t_k, t_{k+N}),$ if $\tilde{x}(t_k) \in \Omega_{\rho_{e,1}}$ $\frac{\partial V(\tilde{x}(t_k))}{\partial x}(f(\tilde{x}(t_k), u(t_k)))$ $\leq \frac{\partial V(\tilde{x}(t_k))}{\partial x} \left(f(\tilde{x}(t_k), h(x(t_k))) \right)$ if $\tilde{x}(t_k) \in \Omega_o / \Omega_{o_{\alpha_1}}$

MOTIVATION FOR HANDLING SIMULTANEOUS ACTUATOR AND SENSOR ATTACKS

• Continuous stirred tank reactor (CSTR) with second-order $A \rightarrow B$ reaction:

$$\frac{dC_A}{dt} = \frac{F}{V}(C_{A0} - C_A) - k_0 e^{\frac{-E}{RT}} C_A^2$$
$$\frac{dT}{dt} = \frac{F}{V}(T_0 - T) + \frac{-\Delta H}{\rho_L C_p} k_0 e^{\frac{-E}{RT}} C_A^2 + \frac{Q}{\rho_L C_p V}$$

- Control objective: Optimize process economics while maintaining the closed-loop state in Ω_{ρ_1}
 - \diamond Economic cost:

$$\int_{t_k}^{t_{k+N}} \left[k_0 e^{-\frac{E}{RT(\tau)}} C_A(\tau)^2\right] d\tau$$

 \diamond Manipulated input constraint

$$0.5 \le C_{A0} \le 7.5 \text{ kmol/m}^3$$

♦ Deviation variables:

$$x_1 = C_A - C_{As}, \quad x_2 = T - T_s$$

 \diamond Process model in input-affine form $\dot{x} = \tilde{f}(x) + gu$

MOTIVATION FOR HANDLING SIMULTANEOUS ACTUATOR AND SENSOR ATTACKS

• Lyapunov-based controller: $h(x) = -1.6x_1 - 0.01x_2$ (M. Heidarinejad, J. Liu, and P. D. Christofides, *SCL*, 2012)

♦ A quadratic Lyapunov function of the form $V_1(x) = x^T P x$ with: $P = \begin{bmatrix} 110.11 & 0 \\ 0 & 0.12 \end{bmatrix}$

♦ Stability region $\rho_1 = 440$ (i.e., $\Omega_{\rho_1} = \{x \in \mathbb{R}^2 : V(x) \le \rho_1\}$)

- $\diamond \ \Omega_{\rho_{e_1}} \subset \Omega_{\rho}, \, \rho_{e_1} = 330$
- LEMPC parameters: $N = 10, \Delta = 0.01$ h
- Process simulated with an integration step size of 10^{-3} h
- The LEMPC receives full state feedback with the full system state $x = [x_1 \ x_2]^T$
- Attack detection policy (initialized at 0.4 h when attack begins): Check if Lyapunov function evaluated at the state measurement decreases over Δ

VARIOUS ATTACK POLICIES

(H. Oyama, D. Messina, K. K. Rangan, and H. Durand, Frontiers in Chemical Engineering, 2022)

- Actuator attack ($u = 0.5 \text{ kmol/m}^3$): Discoverable
- False sensor measurement $(x_1 + 0.5 \text{ kmol/m}^3)$: Not discoverable (no safety issue)
- Combined actuator and sensor attack: Discoverable
- Stealthy actuator and sensor attack (sensor measurements follow trajectory they should have taken): Not discoverable

0.41

♦ State moves closer to safe operating region boundary



PREVENTING SAFETY ISSUES DURING SIMULTANEOUS ATTACKS

- Multiple detector types can be used to aid in cornering an attacker
 - \diamond Examples:
 - Redundant estimators and forcing the decrease of the Lyapunov function across a sampling period
 - ▷ Redundant estimators and state predictions with a redundant control law
 - ♦ Resilient under sufficient conditions
 - Closed-loop state cannot leave a safe operating region in the presence of individual or simultaneous attacks before attack detection
 - Potentially challenging to obtain reasonable control law parameters satisfying resilience theory
- Fundamental notion of cyberattack discoverability:
 - ♦ Whether it is possible to distinguish between a state trajectory coming from attacked sensors and/or actuators and the actual
 - Integrated control and detection policies attempt to use the controller to force differences to show themselves

EXAMPLE OF DISCOVERABILITY-INSPIRED CYBERATTACK DETECTION

(H. Oyama et al., Digital Chemical Engineering, 2023)

- Need a strategy for detecting attacks on sensors that might flag them even with all sensors being compromised
- Set up expectations for measurements that would be "hard" to fake
 - $\diamond\,$ At every sampling time, two control actions are available
 - ♦ Should result in non-overlapping potential sets of measured states
 - $\diamond\,$ One of the two is randomly selected
- Random selection many times in a row
 - ♦ May make it challenging to predict how to not get "caught"



IMPLEMENTING CONTROL ON QUANTUM COMPUTERS

Quantum Computing

- Quantum computing is a technology of recent interest in chemical engineering (D.E Bernal et al., AIChE J., 2022)
- Quantum computers
 exist today of
 different types
- Quantum annealing
 - \diamond Hardware designed to solve certain optimization problems
- Gate-based computers
 - ♦ Considered a path toward "universal" computation



QUANTUM MECHANICS FOR QUANTUM COMPUTING VS. CHEMISTRY

• Reminders from chemistry:

◇ "Time-independent Schrödinger equation" (eigenvalue-eigenvector relationship)

$$\hat{H}(x,t)\psi(x,t) = E\psi(x,t)$$

 \diamond Time-dependent Schrödinger equation

$$\hat{H}(x,t)\psi(x,t) = i\hbar \frac{\partial \psi(x,t)}{\partial t}$$

- $\hat{H}(x,t)$: Hamiltonian (total energy operator)
- Energy *E*
- \hbar : Reduced Planck constant
- $\psi(x,t)$: Wavefunction of the quantum system
 - \diamond Contains information about position of a quantum system
 - \diamond Example: $\psi(x,t)$ is the wavefunction of an electron
 - $\triangleright \psi(x_0, t_0)^* \psi(x_0, t_0) dx$ conveys the probability that the quantum particle will be found in a spatial interval with width dx around x_0 at time t_0 (T. Engel,

Prentice Hall, 2010)

QUANTUM MECHANICS FOR QUANTUM COMPUTING VS. CHEMISTRY

- Wavefunctions are derived from a more fundamental notion of "quantum states"
 - $\diamond "Time-independent Schrödinger equation" (eigenvalue-eigenvector relationship)$

 $\bar{H}(t)\left|\Psi(t)\right\rangle=E\left|\Psi(t)\right\rangle$

 \diamond Time-dependent Schrödinger equation

$$\bar{H}(t) \left| \Psi(t) \right\rangle = i\hbar \frac{\partial \left| \Psi(t) \right\rangle}{\partial t}$$

- $|\Psi(t)\rangle$ is the "quantum state"
 - ♦ "Dirac notation"
- Wavefunctions are derived from the quantum state in a way that makes them particularly good for representing information about position
- Position is continuous
- Gate-based quantum computers generally stay with the binary concept of classical computing
 - \diamond We only want to have 2 possible quantum states for the system
 - \diamond Position will not work for this

QUANTUM MECHANICS FOR QUANTUM COMPUTING VS. CHEMISTRY

- Wavefunctions are derived from a more fundamental notion of "quantum states"
 - $\diamond "Time-independent Schrödinger equation" (eigenvalue-eigenvector relationship)$

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- |Ψ(t)⟩ is the "quantum state"
 ◆ "Dirac notation"
- Wavefunctions are derived from the quantum state in a way that makes them particularly good for representing information about position
- Position is continuous
- Gate-based quantum computers generally stay with the binary concept of classical computing
 - \diamond Wavefunctions are not used in quantum computing
 - \diamond Two possible quantum states: $|0\rangle$ and $|1\rangle$ (regardless of actual implementation)

CONCEPTUALIZING QUANTUM CIRCUITS



- Each unit of a chemical plant changes the state of a process stream
 - Symbols and labeling for process units create meaning for chemical engineers regarding the expected state changes
- Each block ("gate") in a quantum circuit changes the state of a quantum system
 - ♦ Symbols and labeling for the gates create meaning regarding the expected state changes
 - \diamond Example: *H* gate puts a qubit in an equal superposition of two states

QFT-BASED ADDITION

(Ruiz-Perez, L., Garcia-Escartin, J.C., Quantum Information Processing, 2017)



- QFT-based addition: Add two integers a and b (s. Anagolum, Github)
- Binary representations of both numbers are translated to qubit states
- Quantum gates are applied (including those in the inverse QFT, QFT[†]) to obtain final qubit states representative of the bits of the sum

QUANTUM COMPUTING-IMPLEMENTED CONTROL EXAMPLE

Motivation



- Today's quantum computers are noisy
 - ◇ Can cause results of a series of gates to be non-deterministic in practice even if it should be deterministic in theory
- If control was implemented on today's quantum computers, noise could make applied inputs non-deterministic for deterministic process behavior

◇ Raises question of when control could be implemented on quantum computers

• Initial study of these effects: a linear dynamic process, $\dot{x} = x + u$, classically stabilized using the control law u = -2x

QUANTUM COMPUTING-IMPLEMENTED CONTROL EXAMPLE

Noise Model

(Garcia-Escartin, J.C., Chamorro-Posada, P., arXiv, 2011)



- u = -2x is evaluated using a quantum simulator (qasm_simulator) accessed via Qiskit
 - \diamond Use QFT-based addition to compute u=-2x from x+x
- Quantum simulator does not inherently have noise
 - \diamond Required to select a noise model
 - ♦ Evaluated using a controlled Z gate implementation (2 H gates and CNOT gate) as a special case of a controlled phase rotation Z_k

QUANTUM COMPUTING-IMPLEMENTED CONTROL EXAMPLE

Noise Models



• A depolarizing error parameter for qasm_simulator was selected using command for modeling the noise from the 5-qubit quantum device, ibmq_manila, on the qasm_simulator

- ♦ The controlled Z gate was simulated with both the qasm_simulator using this noise model from the device backend and with the depolarizing error parameter set to a fixed value on qasm_simulator
- A depolarizing error parameter of 0.05 was determined to sufficiently approximate the results from the simulations based on ibmq manila

- Comparison between the state trajectories (left) and input trajectories (right) when run with 254 shots for x(0) = 7.4
 - ◇ Classical computer ("Classical system"),
 - $\diamond\,$ Quantum simulator with 254 shots and no noise ("Ideal quantum system")
 - ♦ Quantum simulator with 254 shots and noise ("Noisy quantum system")
- Some deviation is observed between the noisy system and the other two, related to the size (in binary) of the state measurement and number of shots





- Comparison between the state trajectories (left) and input trajectories (right) when run with 1 shot for x(0) = 7.4
 - ◇ Classical computer ("Classical system"),
 - $\diamond\,$ Quantum simulator with 1 shot and no noise ("Ideal quantum system")
 - ◊ Quantum simulator with 1 shot and noise ("Noisy quantum system")
- A significant deviation is observed between the noisy system and the other two, related to the size (in binary) of the state measurement and number of shots





- Comparison between the state trajectories (left) and input trajectories (right) when run with 1 shot for x(0) = 0.74
 - ◇ Classical computer ("Classical system"),
 - $\diamond\,$ Quantum simulator with 1 shot and no noise ("Ideal quantum system")
 - \diamond Quantum simulator with 1 shot and noise ("Noisy quantum system")
- A significant deviation is observed between the noisy system and the other two as a result of the small number of shots





- Comparison between the state trajectories (left) and input trajectories (right) when run with 254 shots for x(0) = 0.74
 - $\diamond\,$ Quantum simulator with 254 shots and noise ("Noisy quantum system")
- No deviation is observed between the noisy system and the other two as a result of the number of shots
- Should we put controllers on quantum computers?
 - ♦ Trying algorithms and evaluating theory to show benefits/limitations





ADVANCED CONTROL AND QUANTUM COMPUTATION

- Rigorous theory for LEMPC makes it attractive for considering the implications of non-deterministic inputs on stability guarantees
 - Initial investigations of closedloop stability of quantum computing-implemented inputs should focus on simple quantum computing algorithms

Table 1: LEMPC solution lookup table

State Measurement	Control Action
0000	1111
0001	1110
0010	1010
	•

- Consider LEMPC solutions in a look-up table

 - $\diamond\,$ Requires quantization of state measurements for LEMPC
 - ♦ Also quantize control actions output by LEMPC

SEARCHING AN LEMPC LOOKUP TABLE VIA MODIFIED GROVER'S SEARCH

- Grover's search algorithm is a quantum computing algorithm for searching an unsorted list (Yanofsky and Mannucci, Cambridge University Press, 2008)
- A modified version of Grover's algorithm could be used to search the LEMPC lookup table
 - ♦ Not efficient for solving this problem
 - Show how non-deterministic inputs can be generated by a quantum computing algorithm tied to LEMPC



- Modified Grover's algorithm implementation strategy:
 - \diamond Use a series of controlled Grover blocks to represent the state/input pairings
 - \diamond Measurements return the "correct" input with probability λ

IMPLICATIONS FOR CLOSED-LOOP STABILITY

- Probability of obtaining the expected control action from Grover's algorithm: λ
- Consider x(t) and $\bar{x}(t) \in \Omega_{\rho_e}$
 - ♦ Control action computed by the LEMPC on a classical computer would maintain $x(t_k)$ and $\bar{x}(t_k)$ in Ω_ρ for $t \in [t_k, t_{k+1})$
 - $\diamond\,$ The modified Grover algorithm would return the same control action as the classical computer with probability λ
 - \diamond Conclusion:
 - $\triangleright \ \mathbf{P}(x(t), \bar{x}(t) \in \Omega_{\rho} \forall \ t \in [t_k, t_{k+1})) \ge \lambda$
- Consider x(t) and $\bar{x}(t) \in \Omega_{\rho}/\Omega_{\rho_e}$
 - ♦ Control action computed by the LEMPC on a classical computer would maintain $x(t_k)$ and $\bar{x}(t_k)$ in Ω_ρ for $t \in [t_k, t_{k+1})$
 - $\diamond\,$ The modified Grover algorithm would return the same control action as the classical computer with probability λ
 - \diamond Conclusion:
 - $\triangleright \mathbf{P}(x(t), \bar{x}(t) \in \Omega_{\rho} \forall t \in [t_k, t_{k+1})) \ge \lambda$

CONCLUSIONS

- Next-generation manufacturing values flexibility and profitability
 - \diamond Facilitated by automation advances such as economic model predictive control
 - ♦ Flexible and profitable systems may not be secure
 - ▷ Attacks on control systems may undermine process safety
- Integrated detection and control policies geared toward nonlinear systems have potential to enable attacks of various types to be detected before causing safety issues
 - $\diamond\,$ Requires sufficient control-theoretic conditions
 - $\diamond\,$ May require at least some sensors to be secure
 - \triangleright Handling attacks after detection likely requires some actuators to be secure
- Fundamental notions of cyberattack-resilience and discoverability for nonlinear systems provide insights into potential future directions for securing controllers
- Quantum computing provides another interesting potential direction for the future of next-generation manufacturing
 - ♦ Control theory and practice require further exploration to determine if benefit exists for quantum computing-implemented control

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